

Differentially Private Oblivious RAM

Sameer Wagh^{*}, Paul Cuff[†], Prateek Mittal^{*}

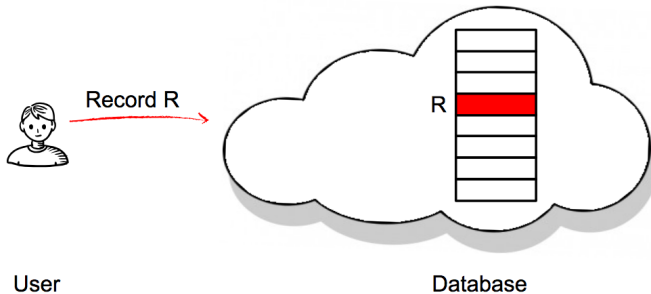
July 24, 2019

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Introduction: Oblivious RAM

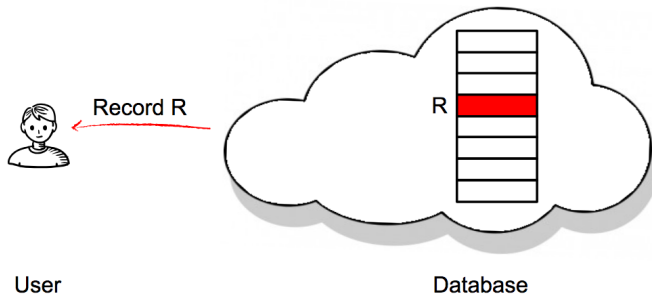
Access data **privately**

from **private** database.



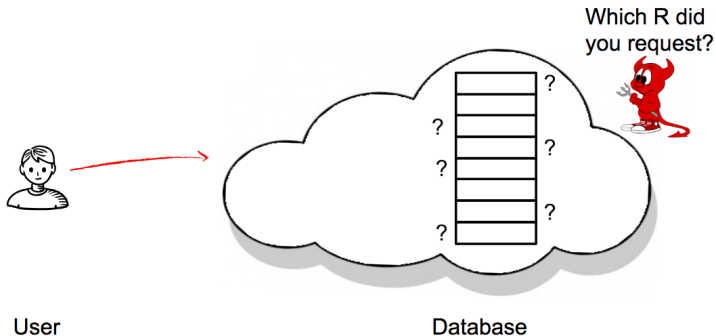
Introduction: Oblivious RAM

User receives record R



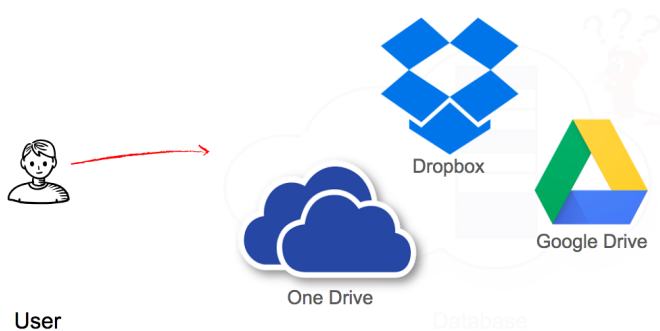
Introduction: Oblivious RAM

Obliviousness: Adversary should not know R



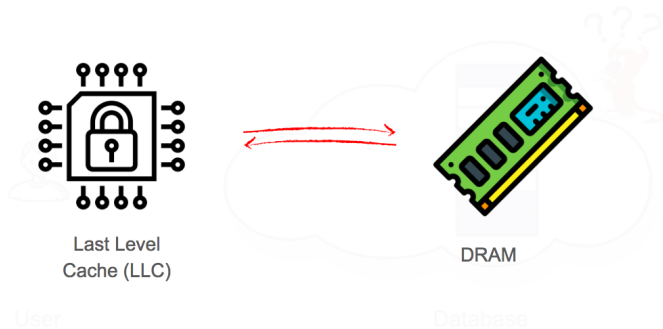
ORAM Application I

Client-server environments

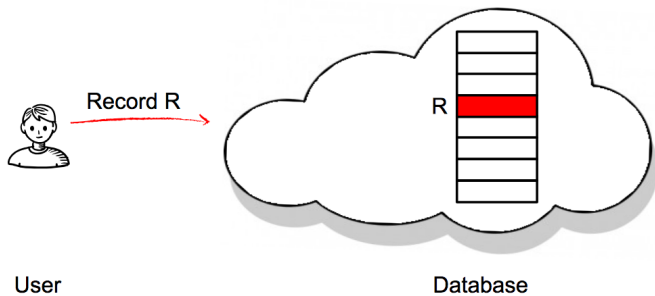


ORAM Application II

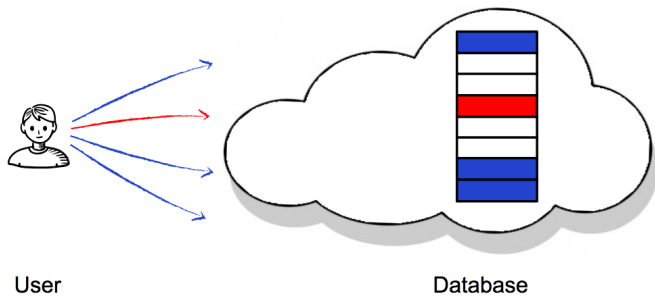
Trusted Execution Environments such as SGX-based enclaves



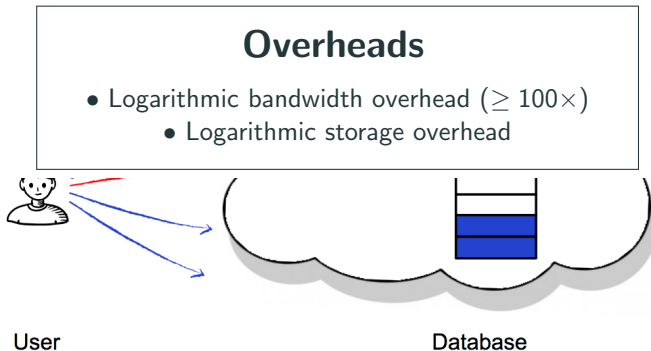
The Problem?



The Problem?



The Problem: Overhead



Key Insight

Can we improve performance by relaxing privacy?

Key Insight: Improve Performance by Relaxing Privacy

- Statistically private ORAM
 - ▶ Better performance at the cost of privacy loss
 - ▶ Challenge: Can we provide rigorous guarantees?

Key Insight: Improve Performance by Relaxing Privacy

- **Statistically private ORAM**
 - ▶ Better performance at the cost of privacy loss
 - ▶ Challenge: Can we provide rigorous guarantees?

- **Efficiency**
 - ▶ Reduce performance overheads – bandwidth, local storage
 - ▶ Achieve privacy proportional to application resources

Differential Privacy

- Formalize Differentially Private ORAM
 - Introduce Root ORAM

Differential Privacy

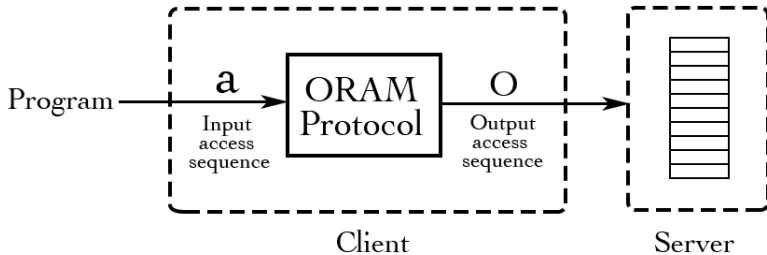
- Formalize Differentially Private ORAM
 - Introduce Root ORAM

Root ORAM

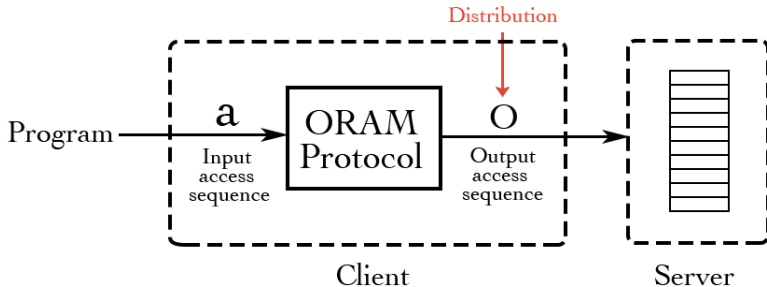
- Theoretical Results
 - Empirical Results
- Private Information Retrieval

Differentially Private Oblivious RAM

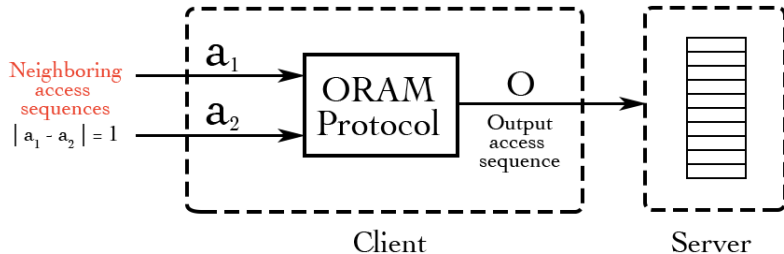
DP-ORAM Intuition



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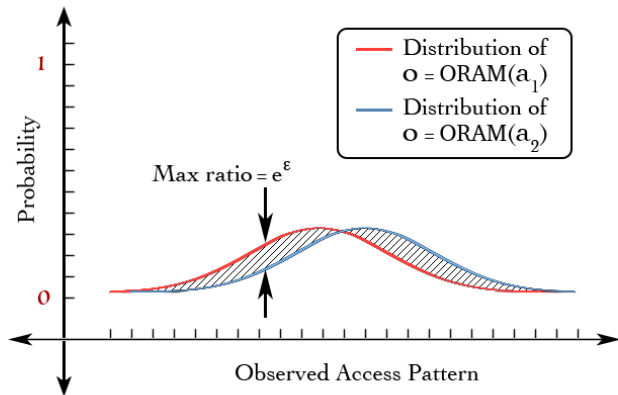


DP-ORAM Intuition



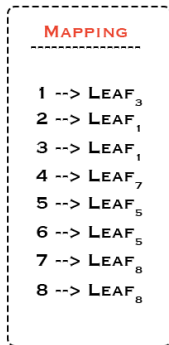
Statistical closeness - Differential Privacy

$$\Pr[\text{ORAM}(\mathbf{a}_1) \in S] \leq e^\epsilon \Pr[\text{ORAM}(\mathbf{a}_2) \in S] + \delta$$

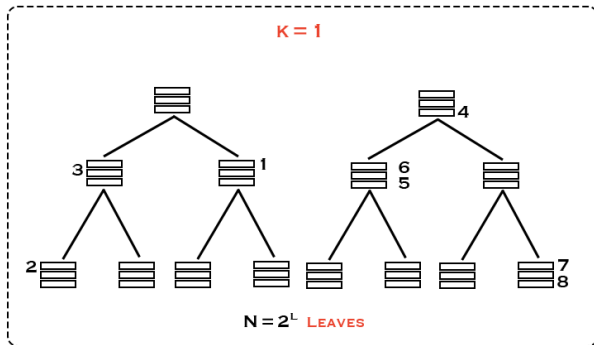


Protocol Construction

Root ORAM: Storage

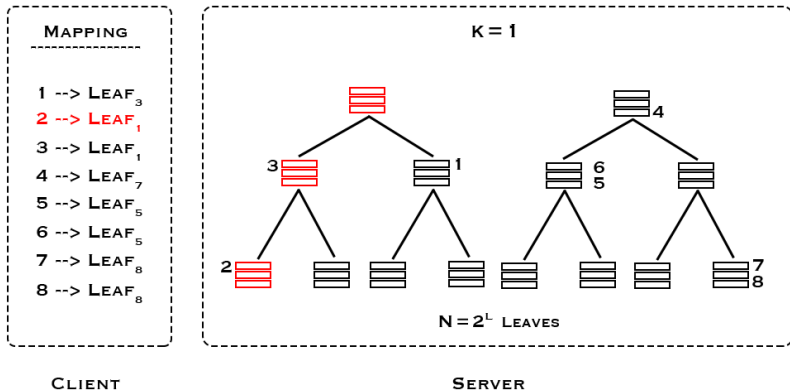


CLIENT

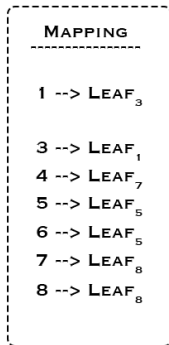


SERVER

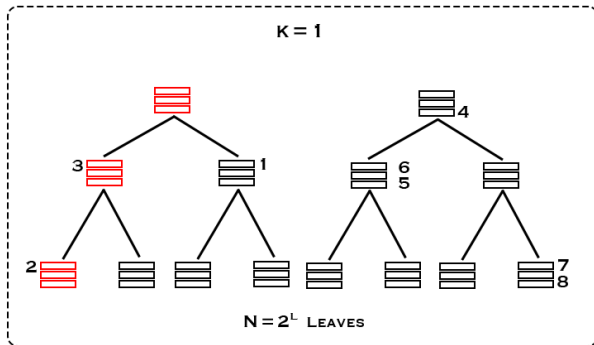
Root ORAM: Invariant



Root ORAM: Updated mapping

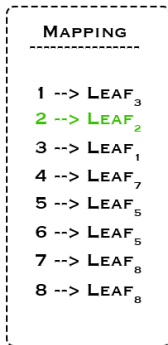


CLIENT

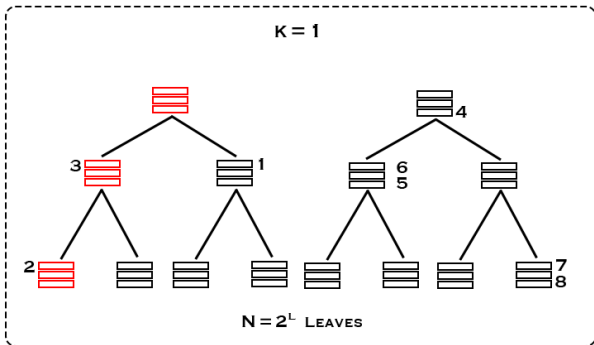


SERVER

Root ORAM: Updated mapping



CLIENT



SERVER

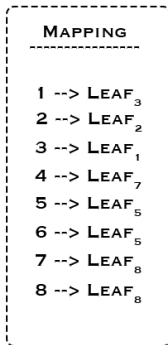
Key Insight

- Uniform mapping \Rightarrow Conventional Security

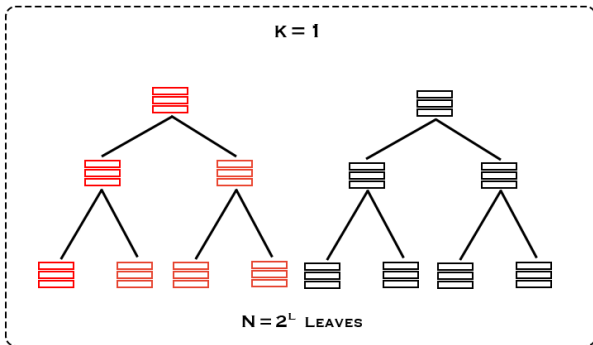
Key Insight

- Uniform mapping \Rightarrow Conventional Security
- Non-uniform mapping \Rightarrow DP-ORAM Security

Root ORAM: Updated mapping



CLIENT

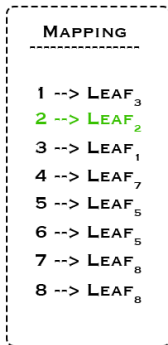


SERVER

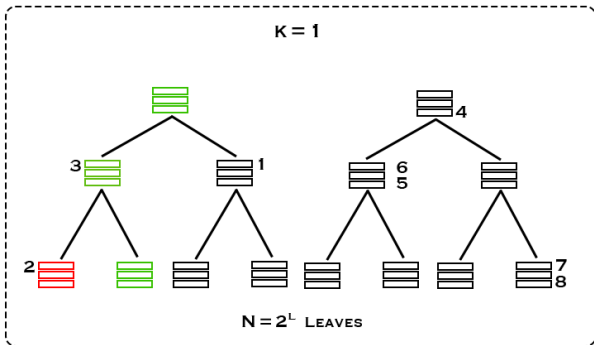
Impact

- Lower average placement \Rightarrow Improved performance
 - Privacy loss

Root ORAM: Write back

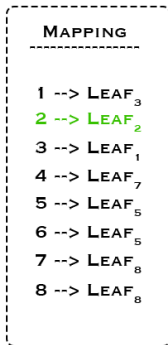


CLIENT

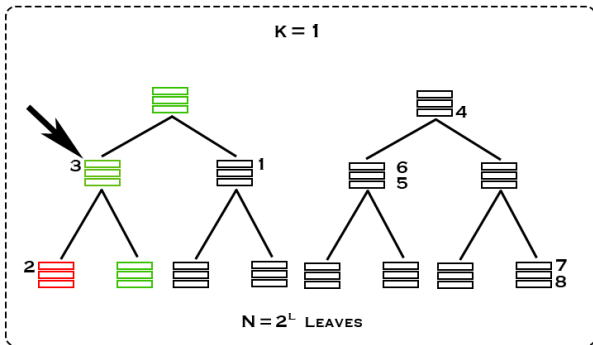


SERVER

Root ORAM: Lowest Common Intersection

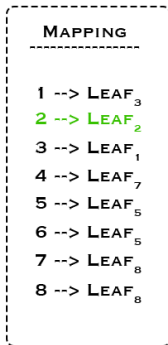


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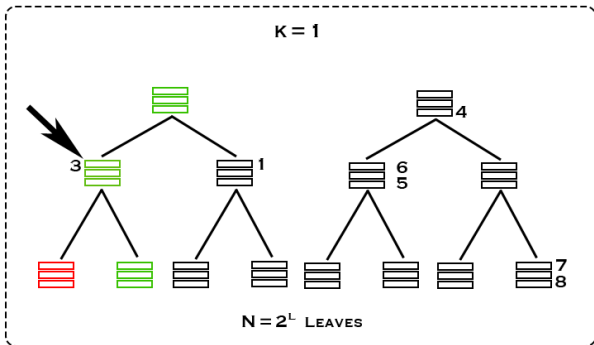


SERVER

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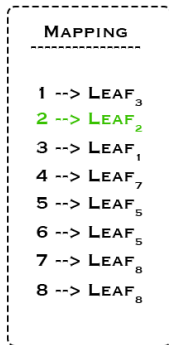


CLIENT

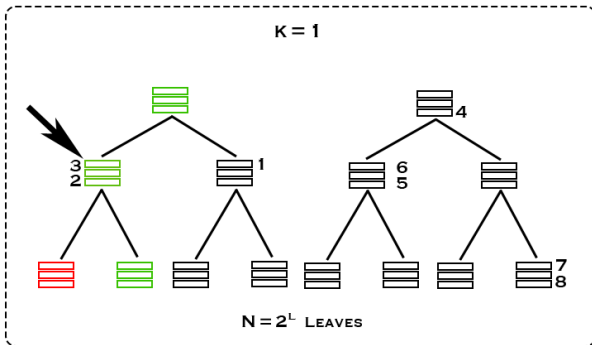


SERVER

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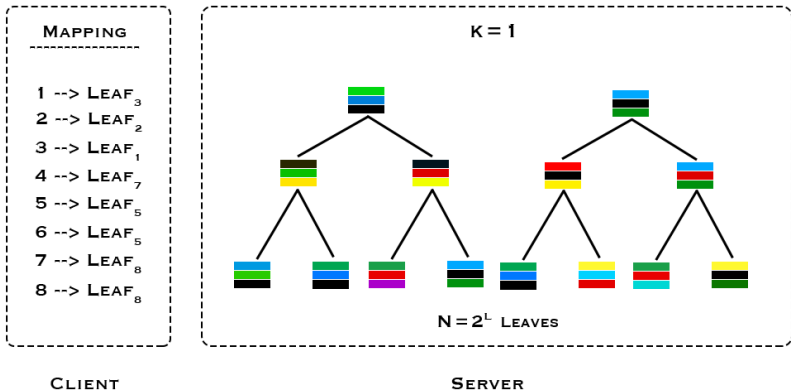


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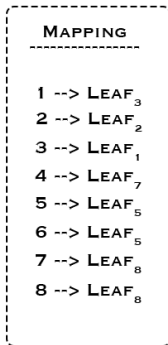


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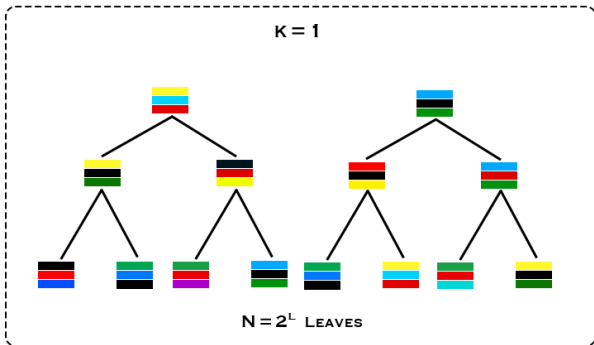
Database view **before** access



Database view **after** access



CLIENT



SERVER

Results

Security Result: Root ORAM is DP-ORAM

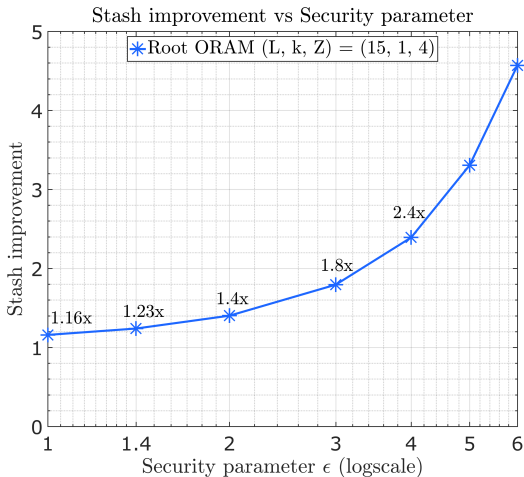
Differentially Private ORAM Protocol

The Root ORAM protocol with parameters k, p is (ϵ, δ) -differentially private for the following choice of ϵ and δ

$$\begin{aligned}\epsilon &= 2 \log \left(\frac{1 + (2^k - 1) \cdot p}{1 - (1 - \delta_{k0})p} \right) \\ \delta &= M \cdot \left(\frac{1 + (2^k - 1) \cdot p}{N} \right)^M\end{aligned}\tag{1}$$

where δ_{k0} is the Kronecker delta, M is the size of the access sequence and $M >$ total stash size.

Performance Improvements



Improvement in stash usage for $(L, k, Z) = (15, 1, 4)$

Key takeaway

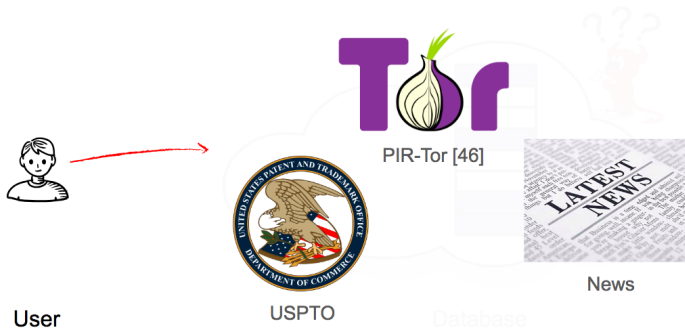
DP-ORAM can enhance performance at the cost of privacy

Application: Private Information Retrieval

Private Information Retrieval (PIR)

Access data **privately**

from **public** database.



[46] Mittal, Prateek, Femi G. Olumofin, Carmela Troncoso, Nikita Borisov, and Ian Goldberg. "PIR-Tor: Scalable Anonymous Communication Using Private Information Retrieval." *In USENIX Security Symposium*, p. 31. 2011.

ORAM based PIR

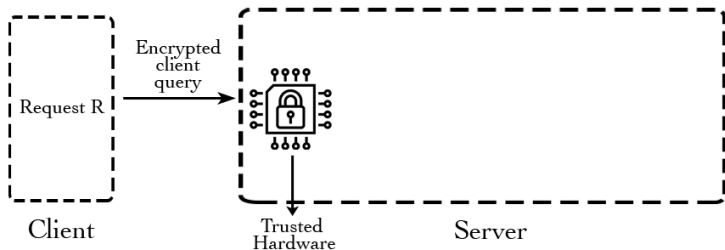
- ORAM has been used previously for PIR [7, 59]

[7] Michael Backes, Aniket Kate, Matteo Maffei, and Kim Pecina. ObliviAd: Provably secure and practical online behavioral advertising. *In IEEE Symposium on Security and Privacy (S&P)*, 2012.

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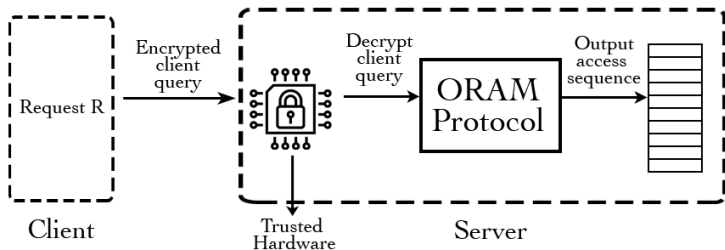


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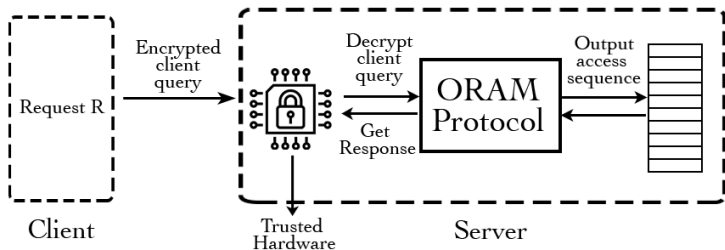


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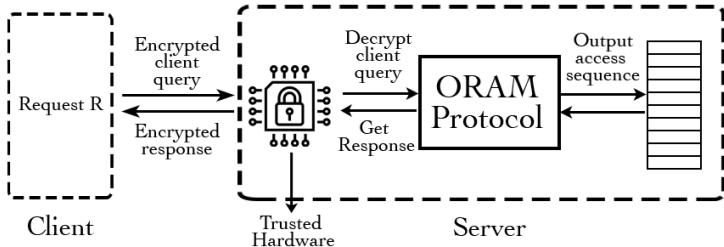


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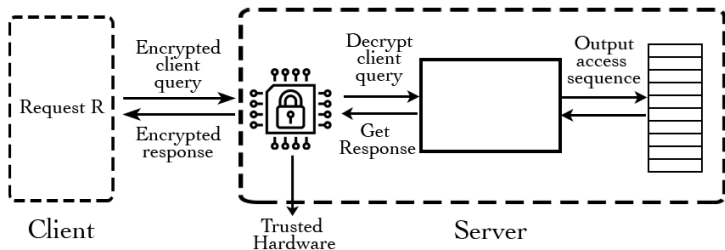


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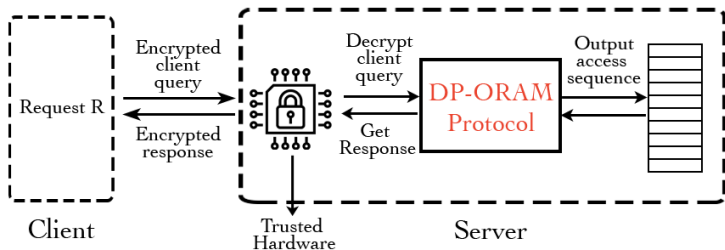


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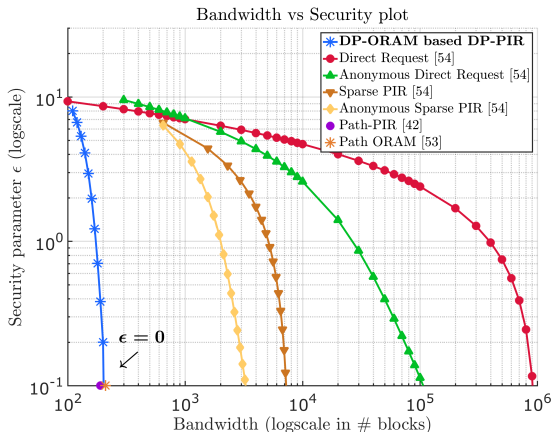
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DP-PIR Bandwidth Comparison



Security-Bandwidth trade-offs for DP-PIR protocols (Toledo *et.al.* [54], Path-PIR [42], and Path ORAM [53]).

DP-PIR Bandwidth Comparison

DP-ORAMs provide significant performance benefits for DP-PIR

Conclusion

Summary

- Formalized Differentially Private ORAMs
- Introduced a family of DP-ORAM protocols
- Analyzed security, performance
- Showcased utility for Private Information Retrieval

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Source code is available at

<https://github.com/inspire-group/Root-ORAM>

Thank you!

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Questions?